

Introduction to Computer Engineering – EECS 203

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Outline

1. Quiz (ungraded)
2. Boolean algebra
3. Homework

Quiz (ungraded)

- What is the relationship between a truth table and a Boolean formula? Is it one/many-to-one/many?
- What is the relationship between a Boolean formula and a combinational circuit? Is it one/many-to-one/many?
- What quirks does the CMOS implementation technology have?
- What is signal restoration or voltage regeneration?
- What practices are unsafe in switch-based design? Why?

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Boolean algebra

- Set of elements, B
- Binary operators, $\{ [\text{AND}, \wedge, *, \cdot], [\text{OR}, \vee, +] \}$
 - We'll prefer \cdot and $+$
 - \cdot frequently omitted
- Unary operator, $[\text{NOT}, ', \bar{}]$

Axioms of Boolean algebra

$$\exists x, y \in B \text{ s.t. } x \neq y$$

closure

$$\forall x, y \in B$$

$$xy \in B$$

$$x + y \in B$$

commutative laws

$$\forall x, y \in B$$

$$xy = yx$$

$$x + y = y + x$$

identities

$$0, 1 \in B, \forall x \in B$$

$$x1 = x$$

$$x + 0 = x$$

Axioms of Boolean algebra

$$\exists x, y \in B \text{ s.t. } x \neq y$$

distributive laws $\forall x, y, z \in B$ $x + (yz) = (x + y)(x + z)$
 $x(y + z) = xy + xz$

complement $x \in B$ $x\bar{x} = 0$
 $x + \bar{x} = 1$

De Morgan's laws

$$\overline{(a + b)} = \bar{a} \bar{b}$$

$$\overline{ab} = \bar{a} + \bar{b}$$

$$\overline{f(x_1, x_2, \dots, x_n, \cdot, +)} = f(\bar{x}_1, \bar{x}_2, \dots, \bar{x}_n, +, \cdot)$$

- Those x s could be functions
- Apply in stages
 - Top-down

De Morgan's laws example

$$\overline{a + bc}$$

$$\bar{a} \cdot \overline{bc}$$

$$\bar{a} \cdot (\bar{b} + \bar{c})$$

Representations of Boolean functions

- Truth table
- Expression using only \cdot , $+$, and $'$
- Symbolic
- Karnaugh map
 - More useful as visualization and optimization tool

Review: *AND*

a	b	a b
0	0	0
0	1	0
1	0	0
1	1	1



$$a \text{ AND } b = a b$$

Will show Karnaugh map later

Review: *OR*

a	b	a + b
0	0	0
0	1	1
1	0	1
1	1	1



$$a \text{ OR } b = a + b$$

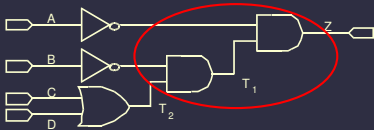
Review: *NOT*

a	\bar{a}
0	1
1	0

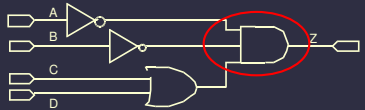


$$\text{NOT } a = \bar{a}$$

Different representations possible



$$Z = ((C + D)\overline{B})\overline{A}$$



$$Z = (C + D)\overline{A}\overline{B}$$

Simplifying logic functions

- Minimize literal count (related to gate count, delay)
- Minimize gate count
- Minimize levels (delay)
- Trade off delay for area
 - Sometimes no real cost

Proving theorems = simplification

Prove $XY + X\bar{Y} = X$

$$\begin{aligned} XY + X\bar{Y} &= X(Y + \bar{Y}) && \text{distributive law} \\ X(Y + \bar{Y}) &= X(1) && \text{complementary law} \\ X(1) &= X && \text{identity law} \end{aligned}$$

Proving theorems = simplification

Prove $X + XY = X$

$$X + XY = X1 + XY \quad \text{identity law}$$

$$X1 + XY = X(1 + Y) \quad \text{distributive law}$$

$$X(1 + Y) = X1 \quad \text{identity law}$$

$$X1 = X \quad \text{identity law}$$

Literals

- Each appearance of a variable (complement) in expression
- Fewer literals usually implies simpler to implement
- E.g., $Z = \overline{A} \overline{B} C + \overline{A} B + \overline{A} B \overline{C} + \overline{B} C$
 - Three variables, ten literals

NANDs and NORs



- Can be implemented in CMOS
 - More on this later
- $X \text{ NAND } Y = \overline{XY}$
- $X \text{ NOR } Y = \overline{X + Y}$
- Do we need inverters?

Summary

- Administrative details
- Finished CMOS and basic gates
- Boolean algebra

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Reading assignment

- M. Morris Mano and Charles R. Kime. *Logic and Computer Design Fundamentals*. Prentice-Hall, NJ, fourth edition, 2008
- Sections 2.4 and 2.5

Next lectures

- Karnaugh maps
- Visual minimization
- We'll also learn the optimal Quine-McCluskey method
 - Optimal two-level minimization is fun!